# Language Processing with Perl and Prolog Chapter 2: Corpus Processing Tools 

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## Corpora

A corpus is a collection of texts (written or spoken) or speech Corpora are balanced from different sources: news, novels, etc.

|  | English | French | German |
| :--- | :--- | :--- | :--- |
| Most frequent words in a collection | the | de | der |
| of contemporary running texts | of | le (article) | die |
|  | to | la (article) | und |
|  | in | et | in |
|  | and | les | des |
| Most frequent words in Genesis | and | et | und |
|  | the | de | die |
|  | of | la | der |
|  | his | à | de |
|  | he | il | er |
|  |  |  | $2 / 39$ |

## Characteristics of Current Corpora

Big: The Bank of English (Collins and U Birmingham) has more than 500 million words
Available in many languages
Easy to collect: The web is the largest corpus ever built and within the reach of a mouse click
Parallel: same text in two languages: English/French (Canadian Hansards),
European parliament (23 languages)
Annotated with part-of-speech or manually parsed (treebanks):

- Characteristics/n of/PREP Current/ADJ Corpora/n
- (NP (NP Characteristics) (PP of (NP Current Corpora)))


## Lexicography

Writing dictionaries
Dictionaries for language learners should be build on real usage

- They're just trying to score brownie points with politicians
- The boss is pleased - that's another brownie point

Bank of English: brownie point (6 occs) brownie points (76 occs) Extensive use of corpora to:

- Find concordances and cite real examples
- Extract collocations and describe frequent pairs of words


## Concordances

A word and its context:

| Language | Concordances |
| :---: | :---: |
| English | s beginning of miracles did Je <br> n they saw the miracles which <br> n can do these miracles that t ain the second miracle that Je <br> e they saw his miracles which |
| French | le premier des miracles que fi i dirent: Quel miracle nous mo om, voyant les miracles qu'il peut faire ces miracles que tu s ne voyez des miracles et des |

## Collocations

Word preferences: Words that occur together

|  | English | French | German |
| :--- | :--- | :--- | :--- |
| You say | Strong tea | Thé fort | Schmales Gesicht |
|  | Powerful computer | Ordinateur puissant | Enge Kleidung |
| You don't | Strong computer | Thé puissant | Schmale Kleidung |
| say | Powerful tea | Ordinateur fort | Enges Gesicht |

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## Word Preferences

| Strong $w$ |  |  | Powerful w |  |  |
| :--- | :--- | :--- | :--- | :--- | :--- |
| strong $w$ | powerful $w$ | $w$ | strong $w$ | powerful $w$ | $w$ |
| 161 | 0 | showing | 1 | 32 | than |
| 175 | 2 | support | 1 | 32 | figure |
| 106 | 0 | defense | 3 | 31 | minority |

## Corpora as Knowledge Sources

Short term:

- Describe usage more accurately
- Assess tools: part-of-speech taggers, parsers.
- Learn statistical/machine learning models for speech recognition, taggers, parsers
- Derive automatically symbolic rules from annotated corpora

Longer term:

- Semantic processing
- Texts are the main repository of human knowledge


## Finite-State Automata

A flexible to tool to search and process text
A FSA accepts and generates strings, here $a c, a b c, a b b c, a b b b c$, abbbbbbbbbbbbc, etc.


## FSA

Mathematically defined by

- $Q$ a finite number of states;
- $\Sigma$ a finite set of symbols or characters: the input alphabet;
- $q_{0}$ a start state,
- $F$ a set of final states $F \subseteq Q$
- $\delta$ a transition function $Q \times \Sigma \rightarrow Q$ where $\delta(q, i)$ returns the state where the automaton moves when it is in state $q$ and consumes the input symbol $i$.


## FSA in Prolog

```
% The start state % The final states
start(q0).
final(q2).
```

transition(q0, a, q1).
transition(q1, b, q1).
transition(q1, c, q2).
accept (Symbols) :-
start(StartState),
accept(Symbols, StartState).
accept([], State) :-
final(State).
accept([Symbol | Symbols], State) :-
transition(State, Symbol, NextState),
accept(Symbols, NextState). accept (Symbols, NextState).

## Regular Expressions

Regexes are equivalent to FSA and generally easier to use Constant regular expressions:

| Pattern | String |
| :--- | :--- |
| regular | A section on regular expressions |
| the | The book of the life |

The automaton above is described by the regex $a b * c$ grep 'ab*c' myFile1 myFile2

## Metacharacters

| Chars | Descriptions | Examples |
| :---: | :---: | :---: |
| * | Matches any number of occurrences of the previous character - zero or more | ac*e matches strings ae, ace, acce, accce, etc. as in "The aerial acceleration alerted the ace pilot" |
| ? | Matches at most one occurrence of the previous character - zero or one | ac?e matches ae and ace as in "The aerial acceleration alerted the ace pilot" |
| + | Matches one or more occurrences of the previous character | ac+e matches ace, acce, accce, etc. as in as in "The aerial acceleration alerted the ace pilot" |

## Metacharacters

| Chars | Descriptions | Examples |
| :--- | :--- | :--- |
| $\{\mathrm{n}\}$ | Matches exactly $n$ occurrences <br> of the previous character | ac\{2\}e matches acce as in <br> "The aerial acceleration alerted <br> the ace pilot" |
| $\{\mathrm{n}\}$, | Matches $n$ or more occurrences <br> of the previous character | ac\{2,\}e matches acce, accce, <br> etc. |
| $\{\mathrm{n}, \mathrm{m}\}$ | Matches from $n$ to $m$ occur- <br> rences of the previous character | ac\{2,4\}e matches acce, <br> acce, and acccce. |

Literal values of metacharacters must be quoted using $\backslash$

## The Dot Metacharacter

The dot . is a metacharacter that matches one occurrence of any character except a new line
a.e matches the strings ale and ace in:

The aerial acceleration alerted the ace pilot
as well as age, ape, are, ate, awe, axe, or aae, aAe, abe, aBe, ale, etc. .* matches any string of characters until we encounter a new line.

## The Longest Match

The previous slide does not tell about the match strategy.
Consider the string $a a b b c$ and the regular expression $\mathrm{a}+\mathrm{b} *$
By default the match engine is greedy: It matches as early and as many characters as possible and the result is aabb
Sometimes a problem. Consider the regular expression <b>.*</b> and the phrase

They match $\langle b\rangle$ as early $\langle/ b\rangle$ and $\langle b\rangle$ as many $</ b\rangle$ characters as they can.

It is possible to use a lazy strategy with the $*$ ? metacharacter instead: <b>.*?</b> and have the result:

They match $\langle b\rangle$ as early $\langle/ b\rangle$ and $\langle b\rangle$ as many $\langle/ b\rangle$ characters as they can.

## Character Classes

[...] matches any character contained in the list.
[^...] matches any character not contained in the list.
[abc] means one occurrence of either $a, b$, or $c$
[ $\sim a b c$ ] means one occurrence of any character that is not an $a, b$, or $c$, [ABCDEFGHIJKLMNOPQRSTUVWXYZ] one upper-case unaccented letter [0123456789] means one digit.
[0123456789]+ . [0123456789]+ matches decimal numbers.
[Cc] omputer [Ss]cience matches Computer Science, computer Science, Computer science, computer science.

## Predefined Character Classes

Expr. Description Example
\d Any digit. Equivalent to [0-9] A\dC matches A0C, A1C, A2C, A3C etc.
$\backslash$ D Any nondigit. Equivalent to [~0-9]
\w Any word character: letter, 1 \w2 matches 1a2, 1A2, 1b2, digit, or underscore. Equivalent 1B2, etc to [a-zA-ZO-9_]
\W Any nonword character. Equivalent to [^\w]
\s Any white space character: space, tabulation, new line, form feed, etc.
$\backslash$ S Any nonwhite space character. Equivalent to [~\s]

## Nonprintable Symbols or Positions

| Char. | Description | Example |
| :---: | :---: | :---: |
| - | Matches the start of a line | - ab*c matches ac, abc, abbc, etc. when they are located at the beginning of a new line |
| \$ | Matches the end of a line | $a b ? c \$$ matches $a c$ and $a b c$ when they are located at the end of a line |
| \b | Matches word boundaries | \babc matches abcd but not dabc <br> bcd $\backslash \mathrm{b}$ matches abcd but not abcde |
| \n | Matches a new line | a \nb matches <br> a <br> b |
| \t | Matches a tabulation | \% |

## Union and Boolean Operators

Union denoted $\mathrm{I}: \mathrm{a} \mid \mathrm{b}$ means either a or b .
Expression $\mathrm{a} \mid \mathrm{bc}$ matches the strings a and bc and $(\mathrm{a} \mid \mathrm{b}) \mathrm{c}$ matches ac and bc,
Order of precedence:
(1) Closure and other repetition operator (highest)
(2) Concatenation, line and word boundaries
(3) Union (lowest)
abc* is the set ab, abc, abcc, abccc, etc.
(abc)* corresponds to abc, abcabc, abcabcabc, etc.

## Perl

Match

```
while ($line = <>) {
    if ($line =~ m/ab+c/) {
        print $line;
        }
}
```


## Substitute

```
while ($line = <>) {
    if ($line =~ m/ab+c/) {
        print "Old: ", $line;
        $line =~ s/ab+c/ABC/g;
        print "New: ", $line;
    }
}
\}
```


## Perl

## Translate

$$
\begin{aligned}
& \text { tr/ABC/abc/ } \\
& \text { \$line }=\sim \operatorname{tr} / \mathrm{A}-\mathrm{Z} / \mathrm{a}-\mathrm{z} / ; \\
& \text { \$line }=\sim \mathrm{tr} / \mathrm{AEIOUaeiou//d;} \\
& \text { \$line }=\sim \operatorname{tr} / A E I O U a e i o u / \$ / \mathrm{cs} ;
\end{aligned}
$$

## Concatenate

```
while ($line = <>) {
    $text .= $line;
}
print $text;
```

References

```
while ($line = <>) {
        while ($line =~ m/\$ *([0-9]+)\.?([0-9]*)/g) {
            print "Dollars: ", $1, " Cents: ", $2,
        }
    }
```

```
\}
```


## Perl

## Predefined variables

\$line = "Tell me, 0 muse, of that ingenious hero who travelled far and wide after he had sacked the famous town of Troy.";
\$line $=^{\sim} \mathrm{m} /, . *, /$;
print \$\&, "\n";
print "Before: ", \$‘, "\n";
print "After: ", \$', "\n";
Arrays
@array $=(1,2,3)$; \#Array containing 1, 2, and 3 print \$array[1]; \#Prints 2

## Concordances in Perl

\# Modified from Doug Cooper
(\$file_name, \$string, \$width) = @ARGV;
open(FILE, "\$file_name")
|l die "Could not open file \$file_name.";
while (\$line = <FILE>) \{
\$text .= \$line;
\}
\$string $=^{\sim}$ s/ /<br>s/g; \# spaces match tabs and new lines \$text $=^{\sim}$ s/\n/ /g; \# new lines are replaced by spaces while (\$text =~ m/(.\{0,\$width\}\$string.\{0,\$width\})/g ) \{ \# matches the pattern with 0..width to the right and left print "\$1\n"; \#\$1 contains the match \}

## Approximate String Matching

A set of edit operations that transforms a source string into a target string: copy, substitution, insertion, deletion, reversal (or transposition). Edits for acress from Kernighan et al. (1990).

| Typo | Correction | Source | Target | Position | Operation |
| :--- | :--- | :---: | :---: | :---: | :--- |
| acress | actress | - | t | 2 | Deletion |
| acress | cress | a | - | 0 | Insertion |
| acress | caress | ac | ca | 0 | Transposition |
| acress | access | r | c | 2 | Substitution |
| acress | across | e | o | 3 | Substitution |
| acress | acres | s | - | 4 | Insertion |
| acress | acres | s | - | 5 | Insertion |
|  |  |  |  |  |  |
|  |  |  |  |  |  |

## Minimum Edit Distance

Edit distances measure the similarity between strings.
We compute the minimum edit distance using a matrix where the value at position $(i, j)$ is defined by the recursive formula:
edit_distance $(i, j)=\min \left(\begin{array}{l}\text { edit__distance }(i-1, j)+\text { del_cost } \\ \text { edit_distance }(i-1, j-1)+\text { subst_cost } \\ \text { edit_distance }(i, j-1)+\text { ins_cost }\end{array}\right)$.
where edit_distance $(i, 0)=i$ and edit_distance $(0, j)=j$.

## Edit Operations



Usually, del_cost = ins_cost =1
subst_cost $=2$ if source $(i) \neq \operatorname{target}(j)$
subst_cost $=0$ if source $(i)=\operatorname{target}(j)$.

## Distance between $a b$ and $c b$



Let us align | a | b | Source |
| :---: | :---: | :---: |
| c | b | Destination |

| b | 2 |  |  |
| :---: | :---: | :---: | :---: |
| c | 1 |  |  |
| Start | 0 | 1 | 2 |
|  | Start | a | b |

## Distance between $a b$ and $c b$



Let us align | a | b | Source |
| :---: | :---: | :---: |
| c | b | Destination |

| b | 2 |  |  |
| :---: | :---: | :---: | :---: |
| c | 1 | 2 |  |
| Start | 0 | 1 | 2 |
|  | Start | a | b |

## Distance between $a b$ and $c b$



Let us align | a | b | Source |
| :---: | :---: | :---: |
| c | b | Destination |

| b | 2 | 3 |  |
| :---: | :---: | :---: | :---: |
| c | 1 | 2 | 3 |
| Start | 0 | 1 | 2 |
|  | Start | a | b |

## Distance between $a b$ and $c b$



Let us align | a | b | Source |
| :---: | :---: | :---: |
| c | b | Destination |

| b | 2 | 3 | $\mathbf{2}$ |
| :---: | :---: | :---: | :---: |
| c | 1 | 2 | 3 |
| Start | 0 | 1 | 2 |
|  | Start | a | b |

## Distance between language and lineage

| e | 7 |  |  |  |  |  |  |  |  |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| g | 6 |  |  |  |  |  |  |  |  |
| a | 5 |  |  |  |  |  |  |  |  |
| e | 4 |  |  |  |  |  |  |  |  |
| n | 3 |  |  |  |  |  |  |  |  |
| i | 2 |  |  |  |  |  |  |  |  |
| l | 1 |  |  |  |  |  |  |  |  |
| Start | 0 | 1 | 2 | 3 | 4 | 5 | 6 | 7 | 8 |
|  | Start | I | a | n | g | u | a | g | e |

## Distance between language and lineage

| e | 7 | 6 | 5 |  |  |  |  |  |  |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| g | 6 | 5 | 4 |  |  |  |  |  |  |
| a | 5 | 4 | 3 |  |  |  |  |  |  |
| e | 4 | 3 | 4 |  |  |  |  |  |  |
| n | 3 | 2 | 3 |  |  |  |  |  |  |
| i | 2 | 1 | 2 | 3 | 4 | 5 | 6 | 7 | 8 |
| l | 1 | 0 | 1 | 2 | 3 | 4 | 5 | 6 | 7 |
| Start | 0 | 1 | 2 | 3 | 4 | 5 | 6 | 7 | 8 |
|  | Start | I | a | n | g | u | a | g | e |

## Distance between language and lineage

| e | 7 | 6 | 5 | 6 | 5 | 6 | 7 | 6 | $\mathbf{5}$ |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| g | 6 | 5 | 4 | 5 | 4 | 5 | 6 | 5 | 6 |
| a | 5 | 4 | 3 | 4 | 5 | 6 | 5 | 6 | 7 |
| e | 4 | 3 | 4 | 3 | 4 | 5 | 6 | 7 | 6 |
| n | 3 | 2 | 3 | 2 | 3 | 4 | 5 | 6 | 7 |
| i | 2 | 1 | 2 | 3 | 4 | 5 | 6 | 7 | 8 |
| l | 1 | 0 | 1 | 2 | 3 | 4 | 5 | 6 | 7 |
| Start | 0 | 1 | 2 | 3 | 4 | 5 | 6 | 7 | 8 |
|  | Start | l | a | n | g | u | a | g | e |

## Perl Code

```
($source, $target) = @ARGV;
$length_s = length($source);
$length_t = length($target);
# Initialize first row and column
for ($i = 0; $i <= $length_s; $i++) {
    $table[$i][0] = $i;
}
for ($j = 0; $j <= $length_t; $j++) {
    $table[0][$j] = $j;
}
# Get the characters. Start index is 0
@source = split(//, $source);
@target = split(//, $target);
```


## Perl Code

```
# Fills the table. Start index of rows and columns is 1
for ($i = 1; $i <= $length_s; $i++) {
    for ($j = 1; $j <= $length_t; $j++) {
    # Is it a copy or a substitution?
        $cost = ($source[$i-1] eq $target[$j-1]) ? 0 : 2;
        # Computes the minimum
        $min = $table[$i-1][$j-1] + $cost;
        if ($min > $table[$i][$j-1] + 1) {
        $min = $table[$i][$j-1] + 1;
        }
        if ($min > $table[$i-1][$j] + 1) {
        $min = $table[$i-1][$j] + 1;
        }
        $table[$i][$j] = $min;
```

    \}
    \}

## Prolog Code

\% edit_operation carries out one edit operation
\% between a source string and a target string.
edit_operation([Char | Source], [Char | Target], Source,
Target, ident, 0).
edit_operation([SChar | Source], [TChar | Target], Source,
Target, sub(SChar,TChar), 2) :-
SChar \= TChar.
edit_operation([SChar | Source], Target, Source, Target, del(SChar), 1).
edit_operation(Source, [TChar | Target], Source, Target, ins(TChar), 1).

## Prolog Code

\% edit_distance(+Source, +Target, -Edits, ?Cost). edit_distance(Source, Target, Edits, Cost) :edit_distance(Source, Target, Edits, 0, Cost).
edit_distance([], [], [], Cost, Cost).
edit_distance(Source, Target, [EditOp | Edits], Cost, FinalCost) :-
edit_operation(Source, Target, NewSource, NewTarget, EditOp, CostOp),
Cost1 is Cost + CostOp, edit_distance(NewSource, NewTarget, Edits, Cost1, FinalCost).

## Distance between language and lineage

First alignment


